

# WARM-START STRATEGIES IN INTERIOR-POINT METHODS FOR LINEAR PROGRAMMING

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**Abstract.** We study the situation in which, having solved a linear program with an interior-point method, we are presented with a new problem instance whose data is slightly perturbed from the original. We describe strategies for recovering a “warm-start” point for the perturbed problem instance from the iterates of the original problem instance. We obtain worst-case estimates of the number of iterations required to converge to a solution of the perturbed instance from the warm-start points, showing that these estimates depend on the size of the perturbation and on the conditioning and other properties of the problem instances.

**1. Introduction.** This paper describes and analyzes warm-start strategies for interior-point methods applied to linear programming (LP) problems. We consider the situation in which one linear program, the “original instance,” has been solved by an interior-point method, and we are then presented with a new problem of the same dimensions, the “perturbed instance,” in which the data is slightly different. Interior-point iterates for the original instance are used to obtain warm-start points for the perturbed instance, so that when an interior-point method is started from this point, it finds the solution in fewer iterations than if no prior information were available. Although our results are theoretical, the strategies proposed here can be applied to practical situations, an aspect that is the subject of ongoing study.

The situation we have outlined arises, for instance, when linearization methods are used to solve nonlinear problems, as in the sequential linear programming algorithm. (One extension of this work that we plan to investigate is to convex quadratic programs, which would be relevant to solution of subproblems in many sequential quadratic programming algorithms.) Our situation is different from the one considered by Gondzio [4], who deals with the case in which the number of unknowns in the primal formulation is increased, and the constraint matrix and cost vector are correspondingly expanded. The latter situation arises in solving subproblems arising from cutting-plane algorithms, for example.

For our analysis, we use the tools developed by Nunez and Freund [5], which in turn are based on the work of Renegar [6, 7, 8, 9] on the conditioning of linear programs and the complexity of algorithms for solving them. We also use standard complexity analysis techniques from the interior-point literature for estimating the number of iterations required to solve a linear program to given accuracy.

We start in Section 2 with an outline of notation and a restatement and slight generalization of the main result from Nunez and Freund [5]. Section 3 outlines the warm-start strategies that we analyze in the paper and describes how our results can be used to obtain reduced complexity estimates for interior-point methods that use the warm starts. In Section 4 we consider a warm-start technique in which a least-squares change is applied to a feasible interior-point iterate for the original instance to make it satisfy the constraints for the perturbed instance. We analyze this technique for central path neighborhoods based on both the Euclidean norm and the  $\infty$  norm, deriving in each case a worst-case estimate for the number of iterations required by

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an interior-point method to converge to an approximate solution of the perturbed instance. In Section 5 we study the technique of applying one iteration of Newton's method to a system of equations that is used to recover a strictly feasible point for the perturbed instance from a feasible iterate for the original instance.

**2. Preliminaries: Conditioning of LPs, Central Path Neighborhoods, Bounds on Feasible Points.** We consider the LP in the following standard form:

$$\min_x c^T x \quad \text{subject to } Ax = b, x \geq 0, \quad (\text{P})$$

where  $A \in R^{m \times n}$ ,  $b \in R^m$ , and  $c \in R^n$  are given and  $x \in R^n$ . The associated dual LP is given by the following:

$$\max_{y,s} b^T y \quad \text{subject to } A^T y + s = c, s \geq 0, \quad (\text{D})$$

where  $y \in R^m$  and  $s \in R^n$ . We borrow the notation of Nunez and Freund [5], denoting by  $d$  the data triplet  $(A, b, c)$  that defines the problems (P) and (D). We define the norm of  $d$  differently from Nunez and Freund, that is, as the maximum of the Euclidean norms of the three data components:

$$(2.1) \quad \|d\| \stackrel{\text{def}}{=} \max(\|A\|_2, \|b\|_2, \|c\|_2).$$

(We will use the norm notation  $\|\cdot\|$  on a vector or matrix to denote the Euclidean norm and the operator norm it induces, respectively, unless explicitly indicated otherwise.)

We use  $\mathcal{F}$  to denote the space of strictly feasible data instances, that is,

$$\mathcal{F} = \{(A, b, c) : \exists x, y, s \text{ with } (x, s) > 0 \text{ such that } Ax = b, A^T y + s = c\}.$$

The complement of  $\mathcal{F}$ , denoted by  $\mathcal{F}^C$ , consists of data instances  $d$  for which either (P) or (D) does not have any strictly feasible solutions. The (shared) boundary of  $\mathcal{F}$  and  $\mathcal{F}^C$  is given by

$$\mathcal{B} = \text{cl}(\mathcal{F}) \cap \text{cl}(\mathcal{F}^C),$$

where  $\text{cl}(\cdot)$  denotes the closure of a set. Since  $(0, 0, 0) \in \mathcal{B}$ , we have that  $\mathcal{B} \neq \emptyset$ . The data instances  $d \in \mathcal{B}$  will be called *ill-posed data instances*, since arbitrary perturbations in the data  $d$  can result in data instances in  $\mathcal{F}$  as well as in  $\mathcal{F}^C$ . The *distance to ill-posedness* is defined as

$$(2.2) \quad \rho(d) = \inf\{\|\Delta d\| : d + \Delta d \in \mathcal{B}\},$$

where we use the norm (2.1) to define  $\|\Delta d\|$ . The condition number of a feasible problem instance  $d$  is defined as

$$(2.3) \quad \mathcal{C}(d) \stackrel{\text{def}}{=} \frac{\|d\|}{\rho(d)}, \quad (\text{with } \mathcal{C}(d) \stackrel{\text{def}}{=} \infty \text{ when } \rho(d) = 0).$$

Since the perturbation  $\Delta d = -d$  certainly has  $d + \Delta d = 0 \in \mathcal{B}$ , we have that  $\rho(d) \leq \|d\|$  and therefore  $\mathcal{C}(d) \geq 1$ . Note, too, that  $\mathcal{C}(d)$  is invariant under a nonzero multiplicative scaling of the data  $d$ , that is,  $\mathcal{C}(\beta d) = \mathcal{C}(d)$  for all  $\beta \neq 0$ .

Robinson [10] and Ashmanov [1] showed that a data instance  $d \in \mathcal{F}$  satisfies  $\rho(d) > 0$  (that is,  $d$  lies in the interior of  $\mathcal{F}$ ) if and only if  $A$  has full row rank. For

such  $d$ , another useful bound on  $\rho(d)$  is provided by the minimum singular value of  $A$ . If we write the singular value decomposition of  $A$  as

$$A = USV^T = \sum_{i=1}^m \sigma_i(A) u_i v_i^T,$$

where  $U$  and  $V$  are orthogonal and  $S = \text{diag}(\sigma_1(A), \sigma_2(A), \dots, \sigma_m(A))$  with  $\sigma_1(A) \geq \sigma_2 \geq \dots \geq \sigma_m(A) > 0$  denoting the singular values of  $A$ , then the perturbation

$$\Delta A = -\sigma_m(A) u_m v_m^T$$

is such that  $A + \Delta A$  is singular, and moreover  $\|\Delta A\| = \sigma_m(A)$  due to the fact that the Euclidean norm of a rank-one matrix satisfies the property

$$(2.4) \quad \|\beta uv^T\|_2 = |\beta| \|u\|_2 \|v\|_2.$$

We conclude that

$$(2.5) \quad \rho(d) \leq \sigma_m(A).$$

It is well known that for such  $d \in \text{int}(\mathcal{F})$ , the system given by

$$(2.6a) \quad Ax = b$$

$$(2.6b) \quad A^T y + s = c$$

$$(2.6c) \quad XSe = \mu e$$

$$(2.6d) \quad (x, s) > 0$$

has a unique solution for every  $\mu > 0$ , where  $e$  denotes the vector of ones in the appropriate dimension and  $X$  and  $S$  are the diagonal matrices formed from the components of  $x$  and  $s$ , respectively. We denote the solutions of (2.6) by  $(x(\mu), y(\mu), s(\mu))$  and use  $\mathcal{P}$  to denote the *central path* traced out by these solutions for  $\mu > 0$ , that is,

$$(2.7) \quad \mathcal{P} \stackrel{\text{def}}{=} \{(x(\mu), y(\mu), s(\mu)) : \mu > 0\}.$$

Throughout this paper, we assume that the original data instance  $d$  lies in  $\mathcal{F}$  and that  $\rho(d) > 0$ . In Sections 4 and 5, we assume further that the original data instance  $d$  has been solved by a feasible path-following interior-point method. Such a method generates a sequence of iterates  $(x^k, y^k, s^k)$  that satisfy the relations (2.6a), (2.6b), and (2.6d) and for which the pairwise products  $x_i^k s_i^k$ ,  $i = 1, 2, \dots, n$ , are not too different from one another, in the sense of remaining within some well-defined “neighborhood” of the central path. The duality measure  $(x^k)^T s^k$  is driven toward zero as  $k \rightarrow \infty$ , and search directions are obtained by applying a modified Newton’s method to the nonlinear system formed by (2.6a), (2.6b), and (2.6c).

We now give some notation for feasible sets and central path neighborhoods associated with the particular problem instance  $d = (A, b, c)$ . Let  $\mathcal{S}$  and  $\mathcal{S}^0$  denote the set of feasible and strictly feasible primal-dual points respectively, that is,

$$\begin{aligned} \mathcal{S} &= \{(x, y, s) : Ax = b, A^T y + s = c, (x, s) \geq 0\}, \\ \mathcal{S}^0 &= \{(x, y, s) \in \mathcal{S} : (x, s) > 0\}. \end{aligned}$$

(Note that  $d \in \mathcal{F}$  if and only if  $\mathcal{S}^0 \neq \emptyset$ .) The central path neighborhoods most commonly used in interior-point methods we refer to as the *narrow* and *wide* neighborhoods. The narrow neighborhood denoted by  $\mathcal{N}_2(\theta)$  is defined as

$$(2.8) \quad \mathcal{N}_2(\theta) = \{(x, y, s) \in \mathcal{S}^0 : \|XSe - (x^T s/n)e\|_2 \leq \theta(x^T s/n)\},$$

for  $\theta \in [0, 1]$ . The wide neighborhood, which is denoted by  $\mathcal{N}_{-\infty}(\gamma)$ , is given by

$$(2.9) \quad \mathcal{N}_{-\infty}(\gamma) = \{(x, y, s) \in \mathcal{S}^0 : x_i s_i \geq \gamma(x^T s/n), \forall i = 1, 2, \dots, n\},$$

where  $u_i$  denotes the  $i$ th component of the vector  $u$  and the parameter  $\gamma$  lies in  $(0, 1]$ .

We typically use a bar to denote the corresponding quantities for the perturbed problem instance  $d + \Delta d$ . That is, we have

$$\begin{aligned} \bar{\mathcal{S}} &= \{(x, y, s) : (A + \Delta A)x = (b + \Delta b), (A + \Delta A)^T y + s = (c + \Delta c), (x, s) \geq 0\}, \\ \bar{\mathcal{S}}^o &= \{(x, y, s) \in \bar{\mathcal{S}} : (x, s) > 0\}, \end{aligned}$$

whereas

$$(2.10a) \quad \bar{\mathcal{N}}_2(\theta) = \{(x, y, s) \in \bar{\mathcal{S}}^o : \|XSe - (x^T s/n)e\|_2 \leq \theta(x^T s/n)\},$$

$$(2.10b) \quad \bar{\mathcal{N}}_{-\infty}(\gamma) = \{(x, y, s) \in \bar{\mathcal{S}}^o : x_i s_i \geq \gamma(x^T s/n), \forall i = 1, 2, \dots, n\}.$$

We associate a value of  $\mu$  with each iterate  $(x, y, s) \in \mathcal{S}$  (or  $\bar{\mathcal{S}}$ ) by setting

$$(2.11) \quad \mu = x^T s/n.$$

We call this  $\mu$  the *duality measure* of the point  $(x, y, s)$ . When  $(x, y, s)$  is feasible, it is easy to show that the duality gap  $c^T x - b^T y$  is equal to  $n\mu$ .

Finally, we state a modified version of Theorem 3.1 from Nunez and Freund [5], which uses our definition (2.1) of the norm of the data instance and takes note of the fact that the proof in [5] continues to hold when we consider strictly feasible points that do not lie exactly on the central path  $\mathcal{P}$ .

**THEOREM 2.1.** *If  $d = (A, b, c) \in \mathcal{F}$  and  $\rho(d) > 0$ , then for any point  $(x, y, s)$  satisfying the conditions*

$$(2.12) \quad Ax = b, \quad A^T y + s = c, \quad (x, s) > 0,$$

*the following bounds are satisfied:*

$$(2.13a) \quad \|x\| \leq \mathcal{C}(d) (\mathcal{C}(d) + \mu n / \|d\|)$$

$$(2.13b) \quad \|y\| \leq \mathcal{C}(d) (\mathcal{C}(d) + \mu n / \|d\|)$$

$$(2.13c) \quad \|s\| \leq 2\|d\|\mathcal{C}(d) (\mathcal{C}(d) + \mu n / \|d\|),$$

*where we have defined  $\mu$  as in (2.11).*

The proof exactly follows the logic of the proof in [5, Theorem 3.1], but differs in many details because of our use of Euclidean norms on the matrices and vectors. For instance, where the original proof defines a perturbation  $\Delta A = -be^T / \|x\|_1$  to obtain an infeasible data instance, we use instead  $\Delta A = -bx^T / \|x\|_2^2$ . We also use the observation (2.4) repeatedly.

**3. Warm Starts and Reduced Complexity.** Before describing specific strategies for warm starts, we preview the nature of our later results and show how they can be used to obtain improved estimates of the complexity of interior-point methods that use these warm starts.

We start by recalling some elements of the complexity analysis of interior-point methods. These methods typically produce iterates  $(x^k, y^k, s^k)$  that lie within a neighborhood such as (2.8) or (2.9) and for which the duality measure  $\mu_k$  (defined as in (2.11) by  $\mu_k = (x^k)^T s^k / n$ ) decreases monotonically with  $k$ , according to a bound of the following form:

$$(3.1) \quad \mu_{k+1} \leq \left(1 - \frac{\delta}{n^\tau}\right) \mu_k,$$

where  $\delta$  and  $\tau$  are positive constants that depend on the algorithm. Typically,  $\tau$  is 0.5, 1, or 2, while  $\delta$  depends on the parameters  $\theta$  or  $\gamma$  that define the neighborhood and various other algorithmic parameters. Given a starting point  $(x^0, y^0, s^0)$  with duality measure  $\mu_0$ , the number of iterations required to satisfy the stopping criterion

$$(3.2) \quad \mu \leq \epsilon \|d\|$$

(for some small positive  $\epsilon$ ) is bounded by

$$(3.3) \quad \frac{\log(\epsilon \|d\|) - \log \mu_0}{\log(1 - \delta/n^\tau)} = \mathcal{O}\left(n^\tau \log \frac{\mu_0}{\|d\| \epsilon}\right).$$

It follows from this bound that, provided we have

$$\frac{\mu_0}{\|d\|} = \mathcal{O}(1/\epsilon^\eta)$$

for some fixed  $\eta > 0$ —which can be guaranteed for small  $\epsilon$  when we apply a cold-start procedure—the number of iterations required to achieve (3.2) is

$$(3.4) \quad \mathcal{O}(n^\tau |\log \epsilon|).$$

Our warm-start strategies aim to find a starting point for the *perturbed* instance that lies inside one of the neighborhoods (2.10), and for which the initial duality measure  $\bar{\mu}_0$  is not too large. By applying (3.3) to the perturbed instance, we see that if  $\bar{\mu}_0 / \|d + \Delta d\|$  is less than 1, the formal complexity of the method will be better than the general estimate (3.4).

Both warm-start strategies that we describe in subsequent sections proceed by taking a point  $(x, y, s)$  from a neighborhood such as (2.8), (2.9) for the original instance and calculating an adjustment  $(\Delta x, \Delta y, \Delta s)$  based on the perturbation  $\Delta d$  to obtain a starting point for the perturbed instance. The strategies are simple; their computational cost is no greater than the cost of one interior-point iteration. They do not succeed in producing a valid starting point unless the point  $(x, y, s)$  from the original problem has a large enough value of  $\mu = x^T s / n$ . That is, we must “back up” along the central path neighborhood until the adjustment  $(\Delta x, \Delta y, \Delta s)$  does not cause some components of  $x$  or  $s$  to become negative. (Indeed, we require a stronger condition to hold: that the adjusted point  $(x + \Delta x, y + \Delta y, s + \Delta s)$  belong to a neighborhood such as those of (2.10).) Since larger perturbations  $\Delta d$  generally lead to larger adjustments  $(\Delta x, \Delta y, \Delta s)$ , the amount by which we must “back up” increases with the size of  $\Delta d$ .

Most of the results in the following sections quantify this observation. They give a lower bound on  $\mu/\|d\|$ —expressed in terms of the size of the components of  $\Delta d$ , the conditioning  $\mathcal{C}(d)$  of the original problem, and other quantities—such that the warm-start strategy applied from a point  $(x, y, s)$  satisfying  $\mu = x^T s/n$  and a neighborhood condition yields a valid starting point for the perturbed problem.

These results can be applied in a practical way when an interior-point approach is used to solve the original instance. Suppose that the iterates  $(x^k, y^k, s^k)$  of this algorithm have been stored and that we restrict the amount by which  $\mu_k$  is decreased on each iteration so that

$$(3.5) \quad \mu_{k+1} \geq \nu \mu_k, \quad \text{for all } k = 0, 1, 2, \dots,$$

for some  $\nu \in (0, 1)$ . Suppose that we denote the lower bound discussed in the preceding paragraph by  $\mu^*/\|d\|$ . Then the best available point for the original instance from which to calculate the warm start is the iterate  $(x^\ell, y^\ell, s^\ell)$ , where  $\ell$  is the largest index for which

$$\mu_\ell \geq \mu^*.$$

Note that because of (3.5) and the choice of  $\ell$ , we have in fact that

$$(3.6) \quad \mu^* \leq \mu_\ell \leq (1/\nu)\mu^*.$$

The warm-start point is then

$$(3.7) \quad (\bar{x}^0, \bar{y}^0, \bar{s}^0) = (x^\ell, y^\ell, s^\ell) + (\Delta x, \Delta y, \Delta s),$$

where  $(\Delta x, \Delta y, \Delta s)$  is the adjustment computed from one of our warm-start strategies. The duality measure corresponding to this point is

$$\bar{\mu}_0 = (\bar{x}^0)^T \bar{s}^0/n = \mu_\ell + (x^\ell)^T \Delta s + (s^\ell)^T \Delta x + \Delta x^T \Delta s.$$

By using the bounds on the components of  $(\Delta x, \Delta y, \Delta s)$  that are obtained during the proofs of each major result in conjunction with the bounds (2.13), we find that  $\bar{\mu}_0$  can be bounded above by some multiple of  $\mu^* + \mu_\ell$ . Because of (3.6), we can deduce in each case that

$$(3.8) \quad \bar{\mu}_0 \leq \kappa \mu^*,$$

for some  $\kappa$  independent of the problem instance  $d$  and the perturbation  $\Delta d$ . We conclude by applying (3.3) to the perturbed instance that the number of iterations required to satisfy the stopping criterion

$$(3.9) \quad \mu \leq \epsilon \|d + \Delta d\|$$

starting from  $(\bar{x}^0, \bar{y}^0, \bar{s}^0)$ , is bounded by

$$(3.10) \quad \mathcal{O} \left( n^\tau \log \frac{\mu^*}{\|d + \Delta d\| \epsilon} \right).$$

Since our assumptions on  $\|\Delta d\|$  usually ensure that

$$(3.11) \quad \|\Delta d\| \leq 0.5 \|d\|,$$

we have that

$$\frac{1}{\|d + \Delta d\|} \leq \frac{1}{\|d\| - \|\Delta d\|} \leq \frac{2}{\|d\|},$$

so that (3.10) can be expressed more conveniently as

$$(3.12) \quad \mathcal{O} \left( n^\tau \log \frac{\mu^*}{\|d\|\epsilon} \right).$$

After some of the results in subsequent sections, we will substitute for  $\tau$  and  $\mu^*$  in (3.12), to express the bound on the number of iterations in terms of the conditioning  $\mathcal{C}(d)$  of the original instance and the size of the perturbation  $\Delta d$ .

Our first warm-start strategy, a least-squares correction, is described in Section 4. The second strategy, a “Newton step correction,” is based on a recent paper by Yıldırım and Todd [12] and is described in Section 5.

**4. Least-Squares Correction.** For much of this section, we restrict our analysis to the changes in  $b$  and  $c$  only; that is, we assume

$$(4.1) \quad \Delta d = (0, \Delta b, \Delta c).$$

Perturbations to  $A$  will be considered in Section 4.3.

Given any primal-dual feasible point  $(x, y, s)$  for the instance  $d$ , the least-squares correction for the perturbation (4.1) is the vector  $(\Delta x, \Delta y, \Delta s)$  obtained from the solutions of the following subproblems:

$$\begin{aligned} \min \|\Delta x\| \quad & \text{s.t. } A(x + \Delta x) = b + \Delta b, \\ \min \|\Delta s\| \quad & \text{s.t. } A^T(y + \Delta y) + (s + \Delta s) = c + \Delta c. \end{aligned}$$

Because  $Ax = b$  and  $A^T y + s = c$ , we can restate these problems as

$$\begin{aligned} \min \|\Delta x\| \quad & \text{s.t. } A\Delta x = \Delta b, \\ \min \|\Delta s\| \quad & \text{s.t. } A^T \Delta y + \Delta s = \Delta c, \end{aligned}$$

which are independent of  $(x, y, s)$ . Given the following QR factorization of  $A^T$ ,

$$(4.2) \quad A^T = \begin{bmatrix} Y & Z \end{bmatrix} \begin{bmatrix} R \\ 0 \end{bmatrix} = YR,$$

where  $\begin{bmatrix} Y & Z \end{bmatrix}$  is orthogonal and  $R$  is upper triangular, we find by simple manipulation of the optimality conditions that the solutions can be written explicitly as

$$(4.3a) \quad \Delta x = YR^{-T}\Delta b,$$

$$(4.3b) \quad \Delta y = R^{-1}Y^T\Delta c,$$

$$(4.3c) \quad \Delta s = (I - YY^T)\Delta c.$$

Observe in particular that

$$(4.4) \quad \Delta x^T \Delta s = 0.$$

Our strategy is as follows: we calculate the correction (4.3) just once, then back-track along the path of iterates  $(x^k, y^k, s^k)$  for the original problem until we find an

index  $k$  such that  $(x^k + \Delta x, s^k + \Delta s) > 0$  and  $(x^k + \Delta x, y^k + \Delta y, s^k + \Delta s)$  lies within either  $\tilde{\mathcal{N}}_2(\theta)$  or  $\tilde{\mathcal{N}}_{-\infty}(\gamma)$ . We hope to be able to satisfy these requirements for some index  $k$  for which the parameter  $\mu_k$  is not too large. In this manner, we hope to obtain a starting point for the perturbed problem for which the initial value of  $\mu$  is not large, so that we can solve the problem using a smaller number of interior-point iterations than if we had started without the benefit of the iterates from the original problem.

Some bounds that we use throughout our analysis follow immediately from (4.3):

$$(4.5) \quad \|\Delta s\| \leq \|\Delta c\|, \quad \|\Delta x\| \leq \frac{\|\Delta b\|}{\sigma_m(A)} \leq \frac{\|\Delta b\|}{\rho(d)},$$

where, as in (2.5),  $\sigma_m(A)$  is the minimum singular value of  $A$ . These bounds follow from the fact that  $I - YY^T$  is an orthogonal projection matrix onto the null space of  $A$  and from the observation that  $R$  has the same singular values as  $A$ . By defining

$$(4.6) \quad \delta_b = \frac{\|\Delta b\|}{\|d\|}, \quad \delta_c = \frac{\|\Delta c\|}{\|d\|},$$

we can rewrite (4.5) as

$$(4.7) \quad \|\Delta s\| \leq \|d\|\delta_c, \quad \|\Delta x\| \leq \mathcal{C}(d)\delta_b.$$

We also define the following quantity, which occurs frequently in the analysis:

$$(4.8) \quad \delta_{bc} = \delta_c + 2\mathcal{C}(d)\delta_b.$$

In the remainder of the paper, we make the mild assumption that

$$(4.9) \quad \delta_b < 1, \quad \delta_c < 1.$$

**4.1. Small Neighborhood.** Suppose that we have iterates for the original problem that satisfy the following property, for some  $\theta_0 \in (0, 1)$ :

$$(4.10) \quad \|XSe - \mu e\|_2 \leq \theta_0 \mu, \quad \text{where } \mu = x^T s / n.$$

That is,  $(x, y, s) \in \mathcal{N}_2(\theta_0)$ . Iterates of a short-step path-following algorithm typically satisfy a condition of this kind. Since  $(x, y, s)$  is a strictly feasible point, its components satisfy the bounds (2.13). Note, too, that we have

$$(4.11) \quad \|XSe - \mu e\| \leq \theta_0 \mu \Rightarrow (1 - \theta_0)\mu \leq x_i s_i \leq (1 + \theta_0)\mu.$$

Our first proposition gives conditions on  $\delta_{bc}$  and  $\mu$  that ensure that the least-squares correction yields a point in the neighborhood  $\tilde{\mathcal{N}}_{-\infty}(\gamma)$ .

**PROPOSITION 4.1.** *Let  $\gamma \in (0, 1 - \theta_0)$  be given, and let  $\xi \in (0, 1 - \gamma - \theta_0)$ . Assume that  $\Delta d$  satisfies*

$$(4.12) \quad \delta_{bc} \leq \frac{1 - \theta_0 - \gamma - \xi}{(n + 1)\mathcal{C}(d)}.$$

*Let  $(x, y, s) \in \mathcal{N}_2(\theta_0)$ , and suppose that  $(\Delta x, \Delta y, \Delta s)$  is the least-squares correction (4.3). Then  $(x + \Delta x, y + \Delta y, s + \Delta s)$  lies in  $\tilde{\mathcal{N}}_{-\infty}(\gamma)$  if*

$$(4.13) \quad \mu \geq \frac{\|d\|}{\xi} 3\mathcal{C}(d)^2 \delta_{bc} \stackrel{\text{def}}{=} \mu_1^*.$$



*Proof.* By using (4.11), (2.13), (4.7), and (4.8), we obtain a lower bound on  $(x_i + \Delta x_i)(s_i + \Delta s_i)$  as follows:

$$\begin{aligned}
& (x_i + \Delta x_i)(s_i + \Delta s_i) \\
&= x_i s_i + x_i \Delta s_i + \Delta x_i s_i + \Delta x_i \Delta s_i \\
&\geq (1 - \theta_0)\mu - \|x\| \|\Delta s\| - \|\Delta x\| \|s\| - \|\Delta x\| \|\Delta s\| \\
&\geq (1 - \theta_0)\mu - \mathcal{C}(d)(\mathcal{C}(d) + \mu n / \|d\|) \|d\| \delta_c \\
&\quad - 2\|d\| \mathcal{C}(d)^2 (\mathcal{C}(d) + \mu n / \|d\|) \delta_b - \|d\| \mathcal{C}(d) \delta_b \delta_c \\
&\geq \mu(1 - \theta_0 - n\mathcal{C}(d)\delta_{bc}) - \mathcal{C}(d)^2 \|d\| \delta_{bc} - \mathcal{C}(d) \|d\| \delta_b \delta_c \\
(4.14) \quad &\geq \mu(1 - \theta_0 - n\mathcal{C}(d)\delta_{bc}) - 2\mathcal{C}(d)^2 \|d\| \delta_{bc}.
\end{aligned}$$

Because of our assumption (4.12), the coefficient of  $\mu$  in (4.14) is positive, so (4.14) represents a positive lower bound on  $(x_i + \Delta x_i)(s_i + \Delta s_i)$  for all  $\mu$  sufficiently large.

For an upper bound on  $(x + \Delta x)^T (s + \Delta s) / n$ , we have from (2.13), (4.7), and the relation (4.4) that

$$\begin{aligned}
& (x + \Delta x)^T (s + \Delta s) / n \\
&\leq \mu + \|\Delta x\| \|s\| / n + \|x\| \|\Delta s\| / n \\
&\leq \mu + 2\mathcal{C}(d)^2 \|d\| \delta_b (\mathcal{C}(d) + \mu n / \|d\|) / n + \mathcal{C}(d) \|d\| \delta_c (\mathcal{C}(d) + \mu n / \|d\|) / n \\
(4.15) \quad &\leq \mu(1 + \mathcal{C}(d)\delta_{bc}) + \mathcal{C}(d)^2 \|d\| \delta_{bc} / n.
\end{aligned}$$

It follows from this bound and (4.14) that a sufficient condition for the conclusion of the proposition to hold is that

$$\mu(1 - \theta_0 - n\mathcal{C}(d)\delta_{bc}) - 2\mathcal{C}(d)^2 \|d\| \delta_{bc} \geq \gamma\mu(1 + \mathcal{C}(d)\delta_{bc}) + \gamma\mathcal{C}(d)^2 \|d\| \delta_{bc} / n,$$

which is equivalent to

$$(4.16) \quad \mu \geq \frac{\|d\| \mathcal{C}(d)^2 \delta_{bc} (2 + \gamma/n)}{1 - \theta_0 - \gamma - \mathcal{C}(d)\delta_{bc}(n + \gamma)},$$

provided that the denominator is positive. Because of the condition (4.12), and using  $\gamma \in (0, 1)$  and  $n \geq 1$ , the denominator is in fact bounded below by the positive quantity  $\xi$ , so the condition (4.16) is implied by (4.13).

Finally, we show that our bounds ensure the positivity of  $x + \Delta x$  and  $s + \Delta s$ . It is easy to show that the right-hand side of (4.14) is also a lower bound on  $(x_i + \alpha \Delta x_i)(s_i + \alpha \Delta s_i)$  for all  $\alpha \in [0, 1]$  and all  $i = 1, 2, \dots, n$ . Because  $\mu$  satisfies (4.16), we have  $(x_i + \alpha \Delta x_i)(s_i + \alpha \Delta s_i) > 0$  for all  $\alpha \in [0, 1]$ . Since we know that  $(x, s) > 0$ , we conclude that  $x_i + \Delta x_i > 0$  and  $s_i + \Delta s_i > 0$  for all  $i$  as well, completing the proof.  $\square$

Next, we seek conditions on  $\delta_{bc}$  and  $\mu$  that ensure that the corrected iterate lies in a narrow central path neighborhood for the perturbed problem.

**PROPOSITION 4.2.** *Let  $\theta > \theta_0$  be given, and let  $\xi \in (0, \theta - \theta_0)$ . Assume that the perturbation  $\Delta d$  satisfies*

$$(4.17) \quad \delta_{bc} \leq \frac{\theta - \theta_0 - \xi}{(2n + 1)\mathcal{C}(d)}.$$

Suppose that  $(x, y, s) \in \mathcal{N}_2(\theta_0)$  for the original problem and that  $(\Delta x, \Delta y, \Delta s)$  is the least-squares correction (4.3). Then,  $(x + \Delta x, y + \Delta y, s + \Delta s)$  will lie in  $\tilde{\mathcal{N}}_2(\theta)$  if

$$(4.18) \quad \mu \geq \frac{\|d\|}{\xi} 4\mathcal{C}(d)^2 \delta_{bc} \stackrel{\text{def}}{=} \mu_2^*.$$

*Proof.* We start by finding a bound on the norm of the vector

$$(4.19) \quad [(x_i + \Delta x_i)(s_i + \Delta s_i)]_{i=1,2,\dots,n} - [(x + \Delta x)^T (s + \Delta s)/n] e.$$

Given two vectors  $y$  and  $z$  in  $R^n$ , we have that

$$(4.20) \quad \left\| [y_i z_i]_{i=1,2,\dots,n} \right\| \leq \|y\| \|z\|, \quad |y^T z| \leq \|y\| \|z\|.$$

By using these elementary inequalities together with (4.4), (4.7), (4.8) and (2.13), we have that the norm of (4.19) is bounded by

$$\begin{aligned} & \left\| [x_i s_i]_{i=1,2,\dots,n} - \mu e \right\| + 2 [\|\Delta x\| \|s\| + \|x\| \|\Delta s\|] + \|\Delta x\| \|\Delta s\| \\ & \leq \theta_0 \mu + 2\mathcal{C}(d) \|d\| \delta_{bc} (\mathcal{C}(d) + n\mu/\|d\|) + \mathcal{C}(d) \|d\| \delta_b \delta_c \\ & \leq [\theta_0 + 2n\mathcal{C}(d) \delta_{bc}] \mu + 3\|d\| \mathcal{C}(d)^2 \delta_{bc}. \end{aligned}$$

Meanwhile, we obtain a lower bound on the duality measure after the correction by using the same set of relations:

$$(4.21) \quad \begin{aligned} (x + \Delta x)^T (s + \Delta s)/n & \geq \mu - [\|\Delta x\| \|s\| + \|x\| \|\Delta s\|]/n \\ & \geq \mu - \mathcal{C}(d) \|d\| \delta_{bc} (\mathcal{C}(d) + n\mu/\|d\|)/n \\ & \geq \mu [1 - \mathcal{C}(d) \delta_{bc}] - \mathcal{C}(d)^2 \|d\| \delta_{bc}/n. \end{aligned}$$

Therefore, a sufficient condition for

$$(x + \Delta x, y + \Delta y, s + \Delta s) \in \tilde{\mathcal{N}}_2(\theta)$$

is that

$$[\theta_0 + 2n\mathcal{C}(d) \delta_{bc}] \mu + 3\|d\| \mathcal{C}(d)^2 \delta_{bc} \leq \theta \mu [1 - \mathcal{C}(d) \delta_{bc}] - \theta \mathcal{C}(d)^2 \|d\| \delta_{bc}/n,$$

which after rearrangement becomes

$$(4.22) \quad \mu [\theta - \theta_0 - 2n\mathcal{C}(d) \delta_{bc} - \theta \mathcal{C}(d) \delta_{bc}] \geq 3\|d\| \mathcal{C}(d)^2 \delta_{bc} + \theta \|d\| \mathcal{C}(d)^2 \delta_{bc}/n.$$

We have from (4.17) that the coefficient of  $\mu$  on the left-hand side of this expression is bounded below by  $\xi$ . By dividing both sides of (4.22) by this expression, and using  $\theta \in (0, 1)$  and  $n \geq 1$ , we find that (4.18) is a sufficient condition for (4.22). A similar argument as in the proof of Proposition 4.1 together with the fact that  $\mu_2^* > \mu_1^*$  ensures positivity of  $(x + \Delta x, s + \Delta s)$ .  $\square$

We now specialize the discussion of Section 3 to show Propositions 4.1 and 4.2 can be used to obtain lower complexity estimates for the interior-point warm-start strategy.

Considering first the case of Proposition 4.1, we have from the standard analysis of a long-step path-following algorithm that constrains its iterates to lie in  $\tilde{\mathcal{N}}_{-\infty}(\gamma)$

(see, for example, Wright [11, Chapter 5]) that the reduction in duality measure at each iteration satisfies (3.1) with

$$\tau = 1, \quad \delta = 2^{\frac{3}{2}} \gamma \frac{1 - \gamma}{1 + \gamma} \min\{\sigma_{\min}(1 - \sigma_{\min}), \sigma_{\max}(1 - \sigma_{\max})\},$$

and  $0 < \sigma_{\min} < \sigma_{\max} < 1$  are the lower and upper bounds on the centering parameter  $\sigma$  at each iteration. Choosing one of the iterates of this algorithm  $(x^\ell, y^\ell, s^\ell)$  in the manner of Section 3, and defining the starting point as in (3.7), we have from (4.15), (4.12), (4.13), and the conditions  $0 < \xi < 1$  and  $n \geq 1$  that

$$\begin{aligned} \bar{\mu}_0 &= (\bar{x}^0)^T \bar{s}^0 / n \\ &\leq \mu_\ell(1 + \mathcal{C}(d)\delta_{bc}) + 2\mathcal{C}(d)^2 \|d\| \delta_{bc} / n \leq \mu_\ell(1 + 1/n) + \mu_1^*(\xi/n) \leq 2\mu_\ell + \mu_1^*. \end{aligned}$$

Now from the property (3.6), it follows that

$$\bar{\mu}_0 \leq (1 + 2/\nu) \mu_1^*.$$

It is easy to verify that (4.12) implies that  $\|\Delta d\| \leq \|d\|/2$ , so that we can use the expression (3.12) to estimate the number of iterations. By substituting  $\tau = 1$  and  $\mu^* = \mu_1^*$  into (3.12), we obtain

$$(4.23) \quad \mathcal{O} \left( n \log \left( \frac{1}{\epsilon} \mathcal{C}(d)^2 \delta_{bc} \right) \right) \quad \text{iterations.}$$

We conclude that if  $\delta_{bc}$  is small in the sense that  $\delta_{bc} \ll \mathcal{C}(d)^{-2}$ , the estimate (4.23) is an improvement on the cold-start complexity estimate (3.4), so it is advantageous to use the warm-start strategy.

Taking now the case of a starting point in the smaller neighborhood of Proposition 4.2, we set  $\theta = 0.4$  and the centering parameter  $\sigma$  to the constant value  $1 - 0.4/n^{1/2}$ . The standard analysis of the short-step path-following algorithm (see, for example, [11, Chapter 4]) then shows that (3.1) holds with

$$\tau = 0.5, \quad \delta = 0.4.$$

By using the procedure outlined in Section 3 to derive the warm-start point, the argument of the preceding paragraph can be applied to obtain the following on the number of iterations:

$$(4.24) \quad \mathcal{O} \left( n^{1/2} \log \left( \frac{1}{\epsilon} \mathcal{C}(d)^2 \delta_{bc} \right) \right).$$

We conclude as before that improved complexity over a cold start is available provided that  $\delta_{bc} \ll \mathcal{C}(d)^{-2}$ .

**4.2. Wide Neighborhood.** We now consider the case in which the iterates for the original problem lie in a wide neighborhood of the central path. To be specific, we suppose that they satisfy  $x_i s_i \geq \gamma_0 \mu$  for some  $\gamma_0 \in (0, 1)$ , that is,  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$ . Note that in this case, we have the following bounds on the pairwise products:

$$(4.25) \quad \gamma_0 \mu \leq x_i s_i \leq (n - (n - 1)\gamma_0) \mu.$$

Similarly to the upper bounds (2.13) on  $\|x\|$  and  $\|s\|$ , we can derive lower bounds on  $x_i$  and  $s_i$  by combining (2.13) with (4.25) and using  $x_i \leq \|x\|$  and  $s_i \leq \|s\|$ :

$$(4.26a) \quad x_i \geq \frac{\gamma_0 \mu}{2\|d\|\mathcal{C}(d)(\mathcal{C}(d) + n\mu/\|d\|)},$$

$$(4.26b) \quad s_i \geq \frac{\gamma_0 \mu}{\mathcal{C}(d)(\mathcal{C}(d) + n\mu/\|d\|)}.$$

These lower bounds will be useful in the later analysis. The following proposition gives a sufficient condition for the least-squares corrected point to be a member of the wide neighborhood for the perturbed problem. The proof uses an argument identical to the proof of Proposition 4.1, with  $\gamma_0$  replacing  $(1 - \theta_0)$ .

**PROPOSITION 4.3.** *Given  $\gamma$  and  $\gamma_0$  such that  $0 < \gamma < \gamma_0 < 1$ , suppose that  $\xi$  is a parameter satisfying  $\xi \in (0, \gamma_0 - \gamma)$ . Assume that  $\Delta d$  satisfies*

$$(4.27) \quad \delta_{bc} \leq \frac{\gamma_0 - \gamma - \xi}{(n+1)\mathcal{C}(d)}.$$

*Suppose that  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$  and denote by  $(\Delta x, \Delta y, \Delta s)$  the least-squares correction (4.3). Then a sufficient condition for*

$$(4.28) \quad (x + \Delta x, y + \Delta y, s + \Delta s) \in \tilde{\mathcal{N}}_{-\infty}(\gamma)$$

*is that*

$$(4.29) \quad \mu \geq \frac{\|d\|}{\xi} 3\mathcal{C}(d)^2 \delta_{bc} \stackrel{\text{def}}{=} \mu_3^*.$$

An argument like the one leading to (4.23) can now be used to show that a long-step path-following method requires

$$(4.30) \quad \mathcal{O} \left( n \log \left( \frac{1}{\epsilon} \mathcal{C}(d)^2 \delta_{bc} \right) \right) \quad \text{iterations}$$

to converge from the warm-start point to a point that satisfies (3.9).

**4.3. Perturbations in  $A$ .** We now allow for perturbations in  $A$  as well as in  $b$  and  $c$ . By doing so, we introduce some complications in the analysis that can be circumvented by imposing an a priori upper bound on the value of  $\mu$  that we are willing to consider. This upper bound is large enough to encompass all values of  $\mu$  of interest from the viewpoint of complexity, in the sense that when  $\mu$  exceeds this bound, the warm-start strategy does not lead to an appreciably improved complexity estimate over the cold-start approach.

For some constant  $\zeta > 1$ , we assume that  $\mu$  satisfies the bound

$$(4.31) \quad \mu \leq \frac{\zeta - 1}{n} \|d\| \mathcal{C}(d) \stackrel{\text{def}}{=} \mu_{\text{up}}$$

so that, for a subexpression that recurs often in the preceding sections, we have

$$\mathcal{C}(d) + n\mu/\|d\| \leq \zeta \mathcal{C}(d).$$

For  $\mu \in [0, \mu_{\text{up}}]$ , we can simplify a number of estimates in the preceding sections, to remove their explicit dependence on  $\mu$ . In particular, the bounds (2.13) on the strictly feasible point  $(x, y, s)$  with  $\mu = x^T s / n$  become

$$(4.32) \quad \|x\| \leq \zeta \mathcal{C}(d)^2, \quad \|y\| \leq \zeta \mathcal{C}(d)^2, \quad \|s\| \leq 2\zeta \|d\| \mathcal{C}(d)^2.$$

Given a perturbation  $\Delta d = (\Delta A, \Delta b, \Delta c)$  with  $\|\Delta d\| < \rho(d)$ , we know that  $A + \Delta A$  has full rank. In particular, for the smallest eigenvalue, we have

$$(4.33) \quad \sigma_m(A + \Delta A) \geq \sigma_m(A) - \|\Delta A\|.$$

To complement the definitions (4.6), we introduce

$$(4.34) \quad \delta_A = \frac{\|\Delta A\|}{\|d\|}.$$

As before, we consider a warm-start strategy obtained by applying least-squares corrections to a given point  $(x, y, s)$  that is strictly feasible for the unperturbed problem. The correction  $\Delta x$  is the solution of the following subproblem:

$$(4.35) \quad \min \|\Delta x\| \quad \text{s.t.} \quad (A + \Delta A)(x + \Delta x) = b + \Delta b,$$

which is given explicitly by

$$(4.36) \quad \Delta x = (A + \Delta A)^T [(A + \Delta A)(A + \Delta A)^T]^{-1} (\Delta b - \Delta A x).$$

By using the QR factorization of  $(A + \Delta A)^T$  as in (4.2) and (4.3), and noting that  $Ax = b$ , we find the following bound on  $\|\Delta x\|$ :

$$(4.37) \quad \|\Delta x\| \leq \frac{\|\Delta b\| + \|\Delta A\|\|x\|}{\sigma_m(A + \Delta A)}.$$

By using (4.33), (2.5), and the definitions (4.6), (4.34), and (2.3), we have

$$\|\Delta x\| \leq \frac{\|\Delta b\| + \|\Delta A\|\|x\|}{\sigma_m(A) - \|\Delta A\|} \leq \frac{\|\Delta b\| + \|\Delta A\|\|x\|}{\rho(d) - \|\Delta A\|} = \frac{\delta_b + \delta_A\|x\|}{1/\mathcal{C}(d) - \delta_A}.$$

In particular, when  $x$  is strictly feasible for the original problem, we have from (4.32) that

$$\|\Delta x\| \leq \mathcal{C}(d) \frac{\delta_b + \zeta \mathcal{C}(d)^2 \delta_A}{1 - \delta_A \mathcal{C}(d)},$$

while if we make the additional simple assumption that

$$(4.38) \quad \delta_A \leq \frac{1}{2\mathcal{C}(d)},$$

we have immediately that

$$(4.39) \quad \|\Delta x\| \leq 2\mathcal{C}(d)\delta_b + 2\zeta\mathcal{C}(d)^3\delta_A.$$

By using (4.38) again, together with (4.9) and the known bounds  $\mathcal{C}(d) \geq 1$  and  $\zeta > 1$ , we obtain

$$(4.40) \quad \|\Delta x\| \leq 2\mathcal{C}(d)\delta_b + 2\zeta\mathcal{C}(d)^3\delta_A \leq 2\mathcal{C}(d) + \zeta\mathcal{C}(d)^2 \leq 3\zeta\mathcal{C}(d)^2.$$

The dual perturbation is the solution of the problem

$$(4.41) \quad \min \|\Delta s\| \quad \text{s.t.} \quad (A + \Delta A)^T (y + \Delta y) + (s + \Delta s) = c + \Delta c.$$

Once again, the minimum norm solution is unique and given by

$$(4.42) \Delta s = \left[ I - (A + \Delta A)^T ((A + \Delta A)(A + \Delta A)^T)^{-1} (A + \Delta A) \right] (\Delta c - \Delta A^T y).$$

Therefore, we have the following upper bound:

$$(4.43) \quad \|\Delta s\| \leq \|\Delta c\| + \|\Delta A\| \|y\|.$$

Using (4.32), we have for  $(x, y, s)$  strictly feasible for the original problem that

$$(4.44) \quad \begin{aligned} \|\Delta s\| &\leq \|\Delta c\| + \|\Delta A\| \zeta \mathcal{C}(d)^2 \\ &\leq \|d\| \delta_c + \zeta \|d\| \mathcal{C}(d)^2 \delta_A. \end{aligned}$$

By using these inequalities, we can prove a result similar to Proposition 4.3.

**PROPOSITION 4.4.** *Suppose we are given  $\gamma$  and  $\gamma_0$  such that  $0 < \gamma < \gamma_0 < 1$ , and a feasible primal-dual point  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$ . Assume further that  $\mu = x^T s / n$  satisfies (4.31) and that the perturbation component  $\Delta A$  satisfies (4.38). For the perturbation  $\Delta d$ , suppose that  $(\Delta x, \Delta y, \Delta s)$  is the least-squares correction obtained from (4.35) and (4.41). We then have*

$$(4.45) \quad (x + \Delta x, y + \Delta y, s + \Delta s) \in \bar{\mathcal{N}}_{-\infty}(\gamma)$$

provided that  $\mu$  satisfies the following lower bound:

$$(4.46) \quad \mu \geq 19 \zeta \mathcal{C}(d)^2 \frac{\|d\|}{\gamma_0 - \gamma} \max(\delta_{bc}, \zeta \mathcal{C}(d)^3 \delta_A) \stackrel{\text{def}}{=} \mu_A^*.$$

*Proof.* By using the upper bounds (4.39) and (4.40) on  $\|\Delta x\|$ , (4.44) on  $\|\Delta s\|$ , and (4.32) on  $\|x\|$  and  $\|s\|$ , we have

$$\begin{aligned} &(x_i + \Delta x_i)(s_i + \Delta s_i) \\ &\geq \gamma_0 \mu - (\|x\| + \|\Delta x\|) \|\Delta s\| - \|s\| \|\Delta x\| \\ &\geq \gamma_0 \mu - [4 \zeta \mathcal{C}(d)^2] [\|d\| \delta_c + \zeta \|d\| \mathcal{C}(d)^2 \delta_A] \\ &\quad - [2 \|d\| \zeta \mathcal{C}(d)^2] [2 \mathcal{C}(d) \delta_b + 2 \zeta \mathcal{C}(d)^3 \delta_A] \\ &\geq \gamma_0 \mu - 4 \|d\| \zeta \mathcal{C}(d)^3 \delta_b - 4 \|d\| \zeta \mathcal{C}(d)^2 \delta_c - 8 \|d\| \zeta^2 \mathcal{C}(d)^5 \delta_A \\ &\geq \gamma_0 \mu - 4 \|d\| \zeta \mathcal{C}(d)^2 \delta_{bc} - 8 \|d\| \zeta^2 \mathcal{C}(d)^5 \delta_A, \end{aligned}$$

where for the last inequality we have used the definition (4.8). By similar logic, and using (4.4), we have for the updated duality measure that

$$\begin{aligned} &(x + \Delta x)^T (s + \Delta s) / n \\ &\leq \mu + \|\Delta x\| \|s\| / n + \|x\| \|\Delta s\| / n \\ &\leq \mu + [2 \mathcal{C}(d) \delta_b + 2 \zeta \mathcal{C}(d)^3 \delta_A] 2 \zeta \|d\| \mathcal{C}(d)^2 / n + \zeta \mathcal{C}(d)^2 [\|d\| \delta_c + \zeta \|d\| \mathcal{C}(d)^2 \delta_A] / n \\ &= \mu + 4 \zeta \mathcal{C}(d)^3 \|d\| \delta_b / n + \zeta \mathcal{C}(d)^2 \|d\| \delta_c / n + 5 \zeta^2 \mathcal{C}(d)^5 \|d\| \delta_A / n \\ &\leq \mu + 2 \zeta \mathcal{C}(d)^2 \|d\| \delta_{bc} / n + 5 \zeta^2 \mathcal{C}(d)^5 \|d\| \delta_A / n. \end{aligned}$$

By comparing these two inequalities in the usual way, and using  $\gamma \in (0, 1)$  and  $n \geq 1$ , we have that a sufficient condition for the conclusion (4.45) to hold is that

$$(4.47) \quad (\gamma_0 - \gamma) \mu \geq 6 \|d\| \zeta \mathcal{C}(d)^2 \delta_{bc} + 13 \|d\| \zeta^2 \mathcal{C}(d)^5 \delta_A.$$

Since from (4.46), we have

$$\begin{aligned}\frac{6}{19}(\gamma_0 - \gamma)\mu &\geq 6\|d\|\zeta\mathcal{C}(d)^2\delta_{bc}, \\ \frac{13}{19}(\gamma_0 - \gamma)\mu &\geq 13\|d\|\zeta^2\mathcal{C}(d)^5\delta_A,\end{aligned}$$

then (4.47) holds, and the proof is complete.  $\square$

By using an argument like the ones leading to (4.23) and (4.30), we deduce that a long-step path-following algorithm that uses the warm start prescribed in Proposition 4.4 requires

$$(4.48) \quad \mathcal{O}\left(n\left[\log\left(\frac{1}{\epsilon}\mathcal{C}(d)^2\delta_{bc}\right) + \log\left(\frac{1}{\epsilon}\mathcal{C}(d)^5\delta_A\right)\right]\right) \quad \text{iterations}$$

to converge to a point that satisfies (3.9).

**5. Newton Step Correction.** In a recent study, Yildirim and Todd [12] analyzed the perturbations in  $b$  and  $c$  in linear and semidefinite programming using interior-point methods. For such perturbations they stated a sufficient condition on the norm of the perturbation, which depends on the current iterate, so that an adjustment to the current point based on applying an iteration of Newton's method to the system (2.6a), (2.6b), (2.6c) yields a feasible iterate for the perturbed problem with a lower duality gap than that of the original iterate. In this section, we augment some of the analysis of [12] with other results, like those of Section 4, to find conditions on the duality gap  $\mu = x^T s/n$  and the perturbation size under which the Newton step yields a warm-start point that yields significantly better complexity than a cold start.

Each iteration of a primal-dual interior-point method involves solving a Newton-like system of linear equations whose coefficient matrix is the Jacobian of the system (2.6a), (2.6b), (2.6c). The general form of these equations is

$$(5.1) \quad \begin{array}{rcl} A\Delta x & = & r_p \\ A^T\Delta y + \Delta s & = & r_d \\ S\Delta x + X\Delta s & = & r_{xs}, \end{array}$$

where typically  $r_p = b - Ax$  and  $r_d = c - A^T y - s$ . The choice of  $r_{xs}$  typically depends on the particular method being applied, but usually represents a Newton or higher-order step toward some “target point”  $(x', y', s')$ , which often lies on the central path  $\mathcal{P}$  defined in (2.7).

In the approach used in Yildirim and Todd [12] and in this section, this Newton-like system is used to correct for perturbations in the data  $(A, b, c)$  rather than to advance to a new primal-dual iterate. The right-hand side quantities are chosen so that that adjustment  $(\Delta x, \Delta y, \Delta s)$  yields a point that is strictly feasible for the perturbed problem, and whose duality gap is no larger than that of the current point  $(x, y, s)$ .

In Section 5.1, we consider the case of perturbations in  $b$  and  $c$  but not in  $A$ . In Section 5.2 we allow perturbations in  $A$  as well.

**5.1. Perturbations in  $b$  and  $c$ .** In our strategy, we assume that

- the current point  $(x, y, s)$  is strictly primal-dual feasible for the original problem;
- the target point  $(x', y', s')$  used to define  $r_{xs}$  is a point that is strictly feasible for the perturbed problem for which  $x'_i s'_i = x_i s_i$  for all  $i = 1, 2, \dots, n$ ;

- the step is a pure Newton step toward  $(x', y', s')$ ; that is,  $r_p = \Delta b$ ,  $r_d = \Delta c$ , and  $r_{xs} = X'S'e - XSe = 0$ .

Note that, in general, the second assumption is not satisfied for an arbitrary current point  $(x, y, s)$  because such a feasible point for the perturbed problem need not exist. However, the Newton's method is still well defined with the above choices of  $r_p$ ,  $r_d$ , and  $r_{xs}$  and that assumption is merely stated for the sake of a complete description of our strategy.

Since  $A$  has full row rank by our assumption of  $\rho(d) > 0$ , we have by substituting our right-hand side in (5.1) and performing block elimination that the solution is given explicitly by

$$(5.2a) \quad \Delta y = (AD^2 A^T)^{-1}(\Delta b + AD^2 \Delta c)$$

$$(5.2b) \quad \Delta s = \Delta c - A^T \Delta y,$$

$$(5.2c) \quad \Delta x = -S^{-1} X \Delta s,$$

where

$$(5.3) \quad D^2 \stackrel{\text{def}}{=} S^{-1} X.$$

Since  $A$  has full row rank and  $D$  is positive diagonal,  $AD^2 A^T$  is invertible.

The following is an extension of the results in Yildirim and Todd [12] to the case of simultaneous perturbations in  $b$  and  $c$ . Note in particular that the Newton step yields a decrease in the duality gap  $x^T s$ .

**PROPOSITION 5.1.** *Assume that  $(x, y, s)$  is a strictly feasible point for  $d$ . Let  $\Delta d = (0, \Delta b, \Delta c)$ . Consider a Newton step  $(\Delta x, \Delta y, \Delta s)$  taken from  $(x, y, s)$  targeting the point  $(x', y', s')$  that is strictly feasible for the perturbed problem and satisfies  $X'S'e = XSe$ , and let*

$$(5.4) \quad (\tilde{x}, \tilde{y}, \tilde{s}) \stackrel{\text{def}}{=} (x, y, s) + (\Delta x, \Delta y, \Delta s).$$

Then if

$$(5.5) \quad \left\| \begin{bmatrix} \Delta c \\ \Delta b \end{bmatrix} \right\|_{\infty} \leq \left\| [S^{-1} (I - A^T (AD^2 A^T)^{-1} AD^2) \quad -S^{-1} A^T (AD^2 A^T)^{-1}] \right\|_{\infty}^{-1},$$

$(\tilde{x}, \tilde{y}, \tilde{s})$  is feasible for the perturbed problem and satisfies

$$(5.6) \quad \tilde{x}^T \tilde{s} \leq x^T s.$$

*Proof.* By rearranging the equation (5.2c) and writing it componentwise, we have

$$(5.7) \quad s_i \Delta x_i + x_i \Delta s_i = 0 \iff \frac{\Delta x_i}{x_i} + \frac{\Delta s_i}{s_i} = 0, \quad i = 1, 2, \dots, n.$$

The next iterate will be feasible if and only if

$$\frac{\Delta x_i}{x_i} \geq -1, \quad \frac{\Delta s_i}{s_i} \geq -1, \quad i = 1, 2, \dots, n.$$



By combining these inequalities with (5.7), we find that feasibility requires that

$$\left| \frac{\Delta x_i}{x_i} \right| \leq 1, \quad \left| \frac{\Delta s_i}{s_i} \right| \leq 1, \quad i = 1, 2, \dots, n,$$

or, equivalently,

$$(5.8) \quad \|S^{-1}\Delta s\|_\infty = \|X^{-1}\Delta x\|_\infty \leq 1.$$

By using (5.2a) and (5.2c), we have

$$\begin{aligned} & \|S^{-1}\Delta s\|_\infty \\ &= \|S^{-1}[\Delta c - A^T \Delta y]\|_\infty \\ (5.9) \quad &= \|S^{-1}[\Delta c - A^T(AD^2 A^T)^{-1}AD^2 \Delta c - A^T(AD^2 A^T)^{-1}\Delta b]\|_\infty \\ &\leq \left\| \begin{bmatrix} S^{-1}(I - A^T(AD^2 A^T)^{-1}AD^2) & -S^{-1}A^T(AD^2 A^T)^{-1} \end{bmatrix} \right\|_\infty \left\| \begin{bmatrix} \Delta c \\ \Delta b \end{bmatrix} \right\|_\infty. \end{aligned}$$

Hence, (5.5) is sufficient to ensure that  $\|S^{-1}\Delta s\|_\infty \leq 1$ .

If we multiply (5.2c) by  $e^T$  from the right, we obtain  $x^T \Delta s + s^T \Delta x = 0$ . Moreover, it follows from (5.7) that  $\Delta x_i$  and  $\Delta s_i$  have opposite signs for each  $i = 1, 2, \dots, n$ , so that  $\Delta x^T \Delta s \leq 0$ . Therefore

$$(x + \Delta x)^T (s + \Delta s) = x^T s + x^T \Delta s + s^T \Delta x + \Delta x^T \Delta s = x^T s + \Delta x^T \Delta s \leq x^T s,$$

proving (5.6).  $\square$

Proposition 5.1 does not provide any insight about the behavior of the expression on the right-hand side of (5.5) as a function of  $\mu$ . To justify our strategy of considering the iterates of the original problem in reverse order, we need to show that the expression in question increases as  $\mu$  corresponding to  $(x, y, s)$  increases, so that we can handle larger perturbations by considering iterates with larger values of  $\mu$ . In the next theorem, we will show that there exists an increasing function  $f(\mu)$  with  $f(0) = 0$  that is a lower bound to the corresponding expression in (5.5) for all values of  $\mu$ . The key to our result is the following bound:

$$(5.10) \quad \chi(H) \stackrel{\text{def}}{=} \sup_{\Sigma \in \mathcal{D}_+} \|\Sigma H^T (H \Sigma H^T)^{-1}\|_\infty < \infty,$$

where  $\mathcal{D}_+$  denotes the set of diagonal matrices in  $R^{n \times n}$  with strictly positive diagonal elements (i.e., positive definite diagonal matrices) and  $\|\cdot\|_\infty$  is the  $\ell_\infty$  matrix norm defined as the maximum of the sums of the absolute values of the entries in each row. This result, by now well known, was apparently first proved by Dikin [2]. For a survey of the background and applications of this and related results, see Forsgren [3].

**THEOREM 5.2.** *Consider points  $(x, y, s)$  in the neighborhood  $\mathcal{N}_\infty(\gamma_0)$  for the original problem, with  $\gamma_0 \in (0, 1)$  and  $\mu = x^T s / n$  as defined in (2.11). Then there exists an increasing function  $f(\mu)$  with  $f(0) = 0$  such that the expression on the right-hand side of (5.5) is bounded below by  $f(\mu)$  for all  $(x, y, s)$  in this neighborhood.*

*Proof.* Let  $(x, y, s)$  be a strictly feasible pair of points for the original problem, which lies in  $\mathcal{N}_\infty(\gamma_0)$  for some  $\gamma_0 \in (0, 1)$ . From (4.26) and (5.10), we have

$$\begin{aligned} & \|S^{-1}A^T(AD^2 A^T)^{-1}\|_\infty = \|S^{-1}D^{-2}D^2 A^T(AD^2 A^T)^{-1}\|_\infty \\ & \leq \|X^{-1}\|_\infty \|D^2 A^T(AD^2 A^T)^{-1}\|_\infty \\ (5.11) \quad & \leq \frac{1}{\mu} \frac{2\|d\|\mathcal{C}(d)}{\gamma_0} (\mathcal{C}(d) + n\mu/\|d\|) \chi(A). \end{aligned}$$

The first inequality is simply the matrix norm inequality. Since  $D^2 = XS^{-1}$ , and  $x$  and  $s$  are strictly feasible,  $D^2$  is a positive definite diagonal matrix, so the bound in (5.10) applies.

Similarly, consider the following:

$$(5.12) \quad \begin{aligned} & \|S^{-1} (I - A^T (AD^2 A^T)^{-1} AD^2)\|_\infty \\ &= \|S^{-1} D^{-1} (I - DA^T (AD^2 A^T)^{-1} AD) D\|_\infty. \end{aligned}$$

Note that  $(I - DA^T (AD^2 A^T)^{-1} AD)$  is a projection matrix onto the nullspace of  $AD$ , therefore, its  $\ell_2$ -norm is bounded by 1. Using the elementary matrix norm inequality  $\|P\|_\infty \leq n^{1/2} \|P\|_2$  for any  $P \in R^{n \times n}$ , we obtain the following sequence of inequalities:

$$(5.13) \quad \begin{aligned} & \|S^{-1} (I - A^T (AD^2 A^T)^{-1} AD^2)\|_\infty \\ &= \|S^{-1} D^{-1} (I - DA^T (AD^2 A^T)^{-1} AD) D\|_\infty \\ &\leq \|X^{-1/2} S^{-1/2}\|_\infty \|I - DA^T (AD^2 A^T)^{-1} AD\|_\infty \|X^{1/2} S^{-1/2}\|_\infty \\ &\leq \max_{i=1,2,\dots,n} \frac{1}{\sqrt{x_i s_i}} n^{1/2} \max_{i=1,2,\dots,n} \sqrt{\frac{x_i}{s_i}} \\ &\leq n^{1/2} \frac{1}{\sqrt{\gamma_0 \mu}} \max_{i=1,2,\dots,n} \frac{x_i}{\sqrt{x_i s_i}} \\ &\leq \frac{1}{\mu} \frac{n^{1/2} \mathcal{C}(d)}{\gamma_0} (\mathcal{C}(d) + n\mu/\|d\|), \end{aligned}$$

where we used  $D^2 = XS^{-1}$ ,  $x_i s_i \geq \gamma_0 \mu$  and (2.13).

If we consider the reciprocal of the right-hand side of the expression (5.5), we obtain

$$(5.14) \quad \begin{aligned} & \| [S^{-1} (I - A^T (AD^2 A^T)^{-1} AD^2) - S^{-1} A^T (AD^2 A^T)^{-1}] \|_\infty \\ &\leq \|S^{-1} (I - A^T (AD^2 A^T)^{-1} AD^2)\|_\infty + \|S^{-1} A^T (AD^2 A^T)^{-1}\|_\infty \\ &\leq \frac{1}{\mu} \frac{2\|d\|\mathcal{C}(d)}{\gamma_0} (\mathcal{C}(d) + n\mu/\|d\|) \chi(A) + \frac{1}{\mu} \frac{n^{1/2} \mathcal{C}(d)}{\gamma_0} (\mathcal{C}(d) + n\mu/\|d\|), \end{aligned}$$

which follows from (5.11) and (5.13). Therefore, (5.14) implies

$$(5.15) \quad \begin{aligned} & \frac{1}{\| [S^{-1} (I - A^T (AD^2 A^T)^{-1} AD^2) - S^{-1} A^T (AD^2 A^T)^{-1}] \|_\infty} \geq \\ & f(\mu) \stackrel{\text{def}}{=} \frac{\gamma_0 \mu}{\mathcal{C}(d) (n^{1/2} + 2\|d\|\chi(A)) [\mathcal{C}(d) + n\mu/\|d\|]}. \end{aligned}$$

It is easy to verify our claims that  $f$  is monotone increasing in  $\mu$  and that  $f(0) = 0$ .  $\square$

Note that Proposition 5.1 guarantees only that the point  $(\tilde{x}, \tilde{y}, \tilde{s})$  is feasible for the perturbed problem. To initiate a feasible path-following interior-point method, we need to impose additional conditions to obtain a strictly feasible point for the perturbed problem that lies in some neighborhood of the central path. For example, in the proof, we imposed only the condition  $(\tilde{x}, \tilde{s}) \geq 0$ . Strict positivity of  $\tilde{x}$  and  $\tilde{s}$  could be ensured by imposing the following condition, for some  $\epsilon \in (0, 1)$ :

$$(5.16) \quad x_i + \Delta x_i \geq \epsilon x_i, \quad s_i + \Delta s_i \geq \epsilon s_i, \quad \forall i = 1, 2, \dots, n.$$

Equivalently, we can replace the necessary and sufficient condition  $\|S^{-1}\Delta s\|_\infty \leq 1$  in (5.8) by the condition  $(\epsilon - 1)e \leq S^{-1}\Delta s \leq (1 - \epsilon)e$ , that is,

$$\|S^{-1}\Delta s\|_\infty \leq 1 - \epsilon,$$

in the proof of Proposition 5.1. With this requirement, we obtain the following bounds:

$$(5.17) \quad \epsilon x_i \leq \tilde{x}_i \leq (2 - \epsilon)x_i, \quad \epsilon s_i \leq \tilde{s}_i \leq (2 - \epsilon)s_i.$$

Note that if  $(\Delta x, \Delta y, \Delta s)$  is the Newton step given by (5.2), then  $\Delta x_i \Delta s_i \leq 0$  for all  $i = 1, 2, \dots, n$ . First, consider the case  $\Delta x_i \geq 0$ , which implies  $\tilde{x}_i \geq x_i$ . We have from (5.17) that

$$(5.18) \quad \tilde{x}_i \tilde{s}_i \geq x_i \tilde{s}_i \geq \epsilon x_i s_i.$$

A similar set of inequalities holds for the case  $\Delta s_i \geq 0$ . Thus, if we define  $\tilde{\mu} = \tilde{x}^T \tilde{s} / n$ , we obtain

$$(5.19) \quad \tilde{\mu} \geq \epsilon \mu.$$

Note that by (5.6), we already have  $\tilde{\mu} \leq \mu$ . With this observation, we can relate the neighborhood in which the original iterate  $(x, y, s)$  lies to the one in which the adjusted point  $(\tilde{x}, \tilde{y}, \tilde{s})$  lies.

**PROPOSITION 5.3.** *Let  $(x, y, s)$  be a strictly feasible point for  $d$ , and suppose that  $\Delta d = (0, \Delta b, \Delta c)$  and  $\epsilon \in (0, 1)$  are given. Consider the Newton step of Proposition 5.1 and the adjusted point  $(\tilde{x}, \tilde{y}, \tilde{s})$  of (5.4). If*

$$(5.20) \quad \left\| \begin{bmatrix} \Delta c \\ \Delta b \end{bmatrix} \right\|_\infty \leq \frac{1 - \epsilon}{\| [S^{-1}(I - A^T(AD^2 A^T)^{-1}AD^2) - S^{-1}A^T(AD^2 A^T)^{-1}] \|_\infty },$$

with  $D$  defined in (5.3), then  $(\tilde{x}, \tilde{y}, \tilde{s})$  is strictly feasible for  $d + \Delta d$  with  $\tilde{\mu} \leq \mu$ . Moreover, if  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$  for the original problem with  $\gamma_0 \in (0, 1)$ , then  $(\tilde{x}, \tilde{y}, \tilde{s})$  satisfies  $(\tilde{x}, \tilde{y}, \tilde{s}) \in \tilde{\mathcal{N}}_{-\infty}(\epsilon\gamma_0)$ .

*Proof.* It suffices to prove the final statement of the theorem. If we assume that  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$ , then using (5.18) and (5.6), we have

$$(5.21) \quad \tilde{x}_i \tilde{s}_i \geq \epsilon x_i s_i \geq \epsilon \gamma_0 \mu \geq \epsilon \gamma_0 \tilde{\mu},$$

which implies that  $(\tilde{x}, \tilde{y}, \tilde{s}) \in \tilde{\mathcal{N}}_{-\infty}(\epsilon\gamma_0)$ , as required.  $\square$

We now have all the tools to be able to prove results like those of Section 4. Suppose that the iterates of the original problem lie in a wide neighborhood with parameter  $\gamma_0$ . For convenience we define

$$(5.22) \quad \|\Delta d\|_\infty \stackrel{\text{def}}{=} \left\| \begin{bmatrix} \Delta b \\ \Delta c \end{bmatrix} \right\|_\infty = \max(\|\Delta b\|_\infty, \|\Delta c\|_\infty).$$

We also define the relative perturbation measure  $\delta_d$  as follows:

$$(5.23) \quad \delta_d \stackrel{\text{def}}{=} \frac{\|\Delta d\|_\infty}{\|d\|}.$$

Note from (4.6) and (4.8) that

$$\delta_d = \max \left( \frac{\|\Delta b\|_\infty}{\|d\|}, \frac{\|\Delta c\|_\infty}{\|d\|} \right) \leq \max(\delta_b, \delta_c) \leq \delta_{bc}.$$

Hence, it is easy to compare results such as Proposition 5.4 below, which obtain a lower bound on  $\mu$  in terms of  $\delta_d$ , to similar results in preceding sections.

Note that Theorem 5.2 provides a lower bound  $f(\mu)$  on the term on the right-hand side of (5.5). Therefore, combining this result with Proposition 5.3, we conclude that a sufficient condition for the perturbation  $\Delta d$  to satisfy (5.20) is that  $\|\Delta d\|_\infty$  is bounded above by the lower bound (5.15) multiplied by  $(1 - \epsilon)$ , that is,

$$\|\Delta d\|_\infty \leq \frac{(1 - \epsilon)\gamma_0\mu}{\mathcal{C}(d)(n^{1/2} + 2\|d\|_\chi(A))(\mathcal{C}(d) + n\mu/\|d\|)},$$

which by rearrangement yields

$$(5.24) \quad \mu \geq \frac{\mathcal{C}(d)^2 \|\Delta d\|_\infty (n^{1/2} + 2\|d\|_\chi(A))}{(1 - \epsilon)\gamma_0 - n\mathcal{C}(d)\|\Delta d\|_\infty (n^{1/2} + 2\|d\|_\chi(A))/\|d\|},$$

provided that the denominator of this expression is positive. To ensure the latter condition, we impose the following bound on  $\delta_d$ :

$$(5.25) \quad \delta_d = \frac{\|\Delta d\|_\infty}{\|d\|} < \frac{(1 - \epsilon)\gamma_0}{n\mathcal{C}(d)(n^{1/2} + 2\|d\|_\chi(A))}.$$

Indeed, when this bound is not satisfied, the perturbation may be so large that the adjusted point  $(\tilde{x}, \tilde{y}, \tilde{s})$  may not be feasible for  $d + \Delta d$  no matter how large we choose  $\mu$  for the original iterate  $(x, y, s)$ .

We now state and prove a result like Proposition 4.3 that gives a sufficient condition on  $\|\Delta d\|_\infty$  and  $\mu$  that ensure that the adjusted point  $(\tilde{x}, \tilde{y}, \tilde{s})$  lies within a wide neighborhood of the central path for the perturbed problem.

**PROPOSITION 5.4.** *Let  $\gamma$  and  $\gamma_0$  be given with  $0 < \gamma < \gamma_0 < 1$ , and suppose that  $\xi$  satisfies  $\xi \in (0, \gamma_0 - \gamma)$ . Assume that  $\delta_d$  satisfies*

$$(5.26) \quad \delta_d \leq \frac{\gamma_0 - \gamma - \xi}{n\mathcal{C}(d)(n^{1/2} + 2\|d\|_\chi(A))}.$$

*Suppose that  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$  for the original problem, and let  $(\tilde{x}, \tilde{y}, \tilde{s})$  be as defined in (5.4). Then if*

$$(5.27) \quad \mu \geq \frac{\|d\|}{\xi} \mathcal{C}(d)^2 \delta_d \left( n^{1/2} + 2\|d\|_\chi(A) \right),$$

*we have  $(\tilde{x}, \tilde{y}, \tilde{s}) \in \tilde{\mathcal{N}}_{-\infty}(\gamma)$ .*

*Proof.* Setting  $\epsilon = \gamma/\gamma_0$ , we note that (5.26) satisfies the condition (5.25), and so the Newton step adjustment yields a strictly feasible point for the perturbed problem. By the argument preceding the proposition, (5.24) gives a sufficient condition for the resulting iterate to lie in  $\tilde{\mathcal{N}}_{-\infty}(\gamma)$  by Proposition 5.3 since  $\gamma = \epsilon\gamma_0$  by the hypothesis. However, (5.26) implies that the denominator of (5.24) is bounded below by  $\xi$ ; hence, (5.24) is implied by (5.27), as required.  $\square$

The usual argument can now be used to show that a long-step path-following method requires

$$(5.28) \quad \mathcal{O} \left( n \log \left( \frac{1}{\epsilon} \mathcal{C}(d)^2 \delta_d \left( n^{1/2} + \|d\|_\chi(A) \right) \right) \right) \quad \text{iterations}$$

to converge from the warm-start point to a point that satisfies (3.9).

**5.2. Perturbations in  $A$ .** In this section, we also allow perturbations in  $A$  (i.e., we let  $\Delta d = (\Delta A, \Delta b, \Delta c)$ ) and propose a Newton step correction strategy to recover warm-start points for the perturbed problem from the iterates of the original problem.

The underlying idea is the same as in Section 5.1. Given a strictly feasible iterate  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$  for the original problem, we apply the Newton's method to recover a feasible point for the perturbed problem by keeping the pairwise products  $x_i s_i$  fixed. As in Section 4.3, we impose an upper bound on  $\mu$  that excludes values of  $\mu$  that are not likely to yield an adjusted starting point with significantly better complexity than a cold-start strategy. In particular, we assume that  $\mu$  satisfies (4.31) for some  $\zeta > 1$ . Let

$$(5.29) \quad \bar{A} \stackrel{\text{def}}{=} A + \Delta A.$$

Given a feasible iterate  $(x, y, s)$  for the original problem, the Newton step correction then is given by the solution to

$$(5.30) \quad \begin{array}{rcl} \bar{A}\Delta x & = & \Delta b - \Delta Ax \\ \bar{A}^T \Delta y + \Delta s & = & \Delta c - \Delta A^T y \\ S\Delta x + X\Delta s & = & 0. \end{array}$$

Under the assumption that  $\bar{A}$  has full row rank, the solution to (5.30) is then given by

$$(5.31a) \quad \Delta y = (\bar{A}D^2 \bar{A}^T)^{-1}(\bar{A}D^2(\Delta c - \Delta A^T y) + \Delta b - \Delta Ax)$$

$$(5.31b) \quad \Delta s = \Delta c - \Delta A^T y - \bar{A}^T \Delta y,$$

$$(5.31c) \quad \Delta x = -S^{-1}X\Delta s,$$

where  $D^2 = S^{-1}X$  as in (5.3).

By a similar argument, a necessary and sufficient condition to have strictly feasible iterates for the perturbed problem is

$$(5.32) \quad \|S^{-1}\Delta s\|_{\infty} \leq 1 - \epsilon, \quad \text{for some } \epsilon \in (0, 1).$$

By Proposition 5.3, the duality gap of the resulting iterate will also be smaller than that of the original iterate. We will modify the analysis in Section 5 to incorporate the perturbation in  $A$  and will refer to the previous analysis without repeating the propositions.

Using (5.31), we get

$$\begin{aligned} S^{-1}\Delta s &= S^{-1}(I - \bar{A}^T(\bar{A}D^2 \bar{A}^T)^{-1}\bar{A}D^2)(\Delta c - \Delta A^T y) \\ &\quad - S^{-1}\bar{A}^T(\bar{A}D^2 \bar{A}^T)^{-1}(\Delta b - \Delta Ax). \end{aligned}$$

Therefore,  $\|S^{-1}\Delta s\|_{\infty}$  is bounded above by

$$\left\| [S^{-1}(I - \bar{A}^T(\bar{A}D^2 \bar{A}^T)^{-1}\bar{A}D^2) - S^{-1}\bar{A}^T(\bar{A}D^2 \bar{A}^T)^{-1}] \right\|_{\infty} \left\| \begin{bmatrix} \Delta c - \Delta A^T y \\ \Delta b - \Delta Ax \end{bmatrix} \right\|_{\infty}.$$

By Theorem 5.2, the first term in this expression is bounded above by  $1/\bar{f}(\mu)$ , where  $\bar{f}(\mu)$  is obtained from  $f(\mu)$  in (5.15) by replacing  $\chi(A)$  by  $\chi(\bar{A})$ . For the second term, we extend the definition (5.22) to account for the perturbations in  $A$  as follows:

$$(5.33) \quad \|\Delta d\|_{\infty} \stackrel{\text{def}}{=} \max(\|\Delta b\|_{\infty}, \|\Delta c\|_{\infty}, \|\Delta A\|_{\infty}, \|\Delta A^T\|_{\infty}),$$

and continue to define  $\delta_d$  as in (5.23). We obtain that

$$\begin{aligned}
& \left\| \begin{bmatrix} \Delta c - \Delta A^T y \\ \Delta b - \Delta A x \end{bmatrix} \right\|_{\infty} \\
& \leq \max\{\|\Delta c\|_{\infty} + \|\Delta A^T\|_{\infty} \|y\|_{\infty}, \|\Delta b\|_{\infty} + \|\Delta A\|_{\infty} \|x\|_{\infty}\}, \\
& \leq \max\{\|\Delta d\|_{\infty} (1 + \|y\|_{\infty}), \|\Delta d\|_{\infty} (1 + \|x\|_{\infty})\}, \\
& \leq \|\Delta d\|_{\infty} (1 + \zeta \mathcal{C}(d)^2), \\
(5.34) \quad & \leq 2\|\Delta d\|_{\infty} \zeta \mathcal{C}(d)^2,
\end{aligned}$$

where we used (5.33), (4.32),  $\zeta > 1$  and  $\mathcal{C}(d) \geq 1$  to derive the inequalities. By combining the two upper bounds we obtain

$$(5.35) \quad \|S^{-1} \Delta s\|_{\infty} \leq \frac{1}{\mu} \frac{1}{\gamma_0} 2\zeta \mathcal{C}(d)^3 \left( n^{1/2} + 2\|d\|_{\chi(\bar{A})} \right) (\mathcal{C}(d) + n\mu/\|d\|) \|\Delta d\|_{\infty}.$$

Therefore, a sufficient condition to ensure (5.32) is obtained by requiring the upper bound in (5.35) to be less than  $1 - \epsilon$ . Rearranging the resulting inequality yields a lower bound on  $\mu$ :

$$(5.36) \quad \mu \geq \frac{2\zeta \mathcal{C}(d)^4 (n^{1/2} + 2\|d\|_{\chi(\bar{A})}) \|\Delta d\|_{\infty}}{\gamma_0 (1 - \epsilon) - 2\zeta n \mathcal{C}(d)^3 (n^{1/2} + 2\|d\|_{\chi(\bar{A})}) \|\Delta d\|_{\infty} / \|d\|},$$

provided that the denominator is positive, which is ensured by the condition

$$(5.37) \quad \delta_d = \frac{\|\Delta d\|_{\infty}}{\|d\|} < \frac{\gamma_0 (1 - \epsilon)}{2\zeta n \mathcal{C}(d)^3 (n^{1/2} + 2\|d\|_{\chi(\bar{A})})}.$$

The proof of the following result is similar to that of Proposition 5.4.

**PROPOSITION 5.5.** *Let  $\gamma$  and  $\gamma_0$  be given with  $0 < \gamma < \gamma_0 < 1$ , and suppose that  $\xi$  satisfies  $\xi \in (0, \gamma_0 - \gamma)$ . Assume that  $\Delta d$  satisfies*

$$(5.38) \quad \delta_d \leq \frac{\gamma_0 - \gamma - \xi}{2\zeta n \mathcal{C}(d)^3 (n^{1/2} + 2\|d\|_{\chi(\bar{A})})}.$$

*Suppose that  $(x, y, s) \in \mathcal{N}_{-\infty}(\gamma_0)$  and that  $(\tilde{x}, \tilde{y}, \tilde{s})$  is the adjusted point defined in (5.4). Then we have  $(\tilde{x}, \tilde{y}, \tilde{s}) \in \mathcal{N}_{-\infty}(\gamma)$  provided that*

$$(5.39) \quad \mu \geq \frac{\|d\|}{\xi} 2\zeta \mathcal{C}(d)^4 \delta_d \left( n^{1/2} + 2\|d\|_{\chi(\bar{A})} \right).$$

The usual argument can be used again to show that a long-step path-following method requires

$$(5.40) \quad \mathcal{O} \left( n \log \left( \frac{1}{\epsilon} \mathcal{C}(d)^4 \delta_d \left( n^{1/2} + \|d\|_{\chi(\bar{A})} \right) \right) \right) \quad \text{iterations}$$

to converge from the warm-start point to a point that satisfies (3.9).

**6. Conclusions.** We have described two schemes by which the iterates of an interior-point method applied to an LP instance can be adjusted to obtain starting points for a perturbed instance. We have derived worst-case estimates for the number of iterations required to obtain convergence from these warm starting points. These estimates depend chiefly on the size of the perturbation, on the conditioning of the original problem instance, and on a key property of the constraint matrix.

In future work, we plan to extend the techniques to infeasible-interior-point methods, and perform computational experiments to determine the practical usefulness of these techniques. We will also investigate extensions to wider classes of problems, such as convex quadratic programs and linear complementarity problems.

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